

# Survey on approximation methods for discrete scan statistics: a software illustration

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## The $d$ -dimensional discrete scan statistics

Let  $T_1, T_2, \dots, T_d$  be positive integers, with  $d \geq 1$

- The rectangular region,  $\mathcal{R}_d = [0, T_1] \times [0, T_2] \times \dots \times [0, T_d]$
- The r.v.'s  $X_{s_1, s_2, \dots, s_d}$ ,  $1 \leq s_j \leq T_j$ ,  $j \in \{1, 2, \dots, d\}$

Let  $2 \leq m_j \leq T_j$ ,  $1 \leq j \leq d$ , be positive integers

- Define for  $1 \leq i_l \leq T_l - m_l + 1$ ,  $1 \leq l \leq d$ ,

$$Y_{i_1, i_2, \dots, i_d} = \sum_{s_1=i_1}^{i_1+m_1-1} \sum_{s_2=i_2}^{i_2+m_2-1} \dots \sum_{s_d=i_d}^{i_d+m_d-1} X_{s_1, s_2, \dots, s_d}$$

- The  $d$ -dimensional discrete scan statistic,

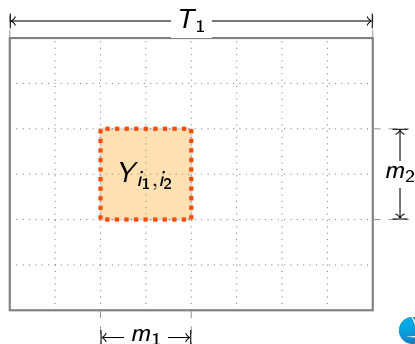
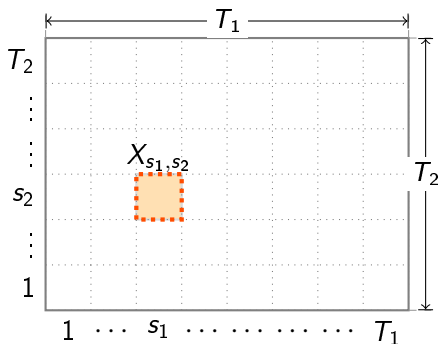
$$S_{\mathbf{m}}(\mathbf{T}) = \max_{\substack{1 \leq j_l \leq T_l - m_l + 1 \\ j \in \{1, 2, \dots, d\}}} Y_{i_1, i_2, \dots, i_d}$$

with  $\mathbf{m} = (m_1, m_2, \dots, m_d)$  and  $\mathbf{T} = (T_1, T_2, \dots, T_d)$

# Example: two dimensional scan statistics ( $d = 2$ )

We have for  $d = 2$

$$Y_{i_1, i_2} = \sum_{s_1=i_1}^{i_1+m_1-1} \sum_{s_2=i_2}^{i_2+m_2-1} X_{s_1, s_2}, \quad S_{m_1, m_2}(T_1, T_2) = \max_{\substack{1 \leq i_1 \leq T_1 - m_1 + 1 \\ 1 \leq i_2 \leq T_2 - m_2 + 1}} Y_{i_1, i_2}$$



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# Problem

## Goal

Find a good estimate for the distribution of  $d$ -dimensional discrete scan statistic

$$Q_m(\mathbf{T}) = \mathbb{P}(S_m(\mathbf{T}) \leq \tau)$$

The distribution of  $S_m(\mathbf{T})$  is used for testing the null hypotheses of randomness against the alternative hypothesis of clustering.

## Example: Bernoulli model

$H_0$ : The r.v.'s  $X_{s_1, s_2, \dots, s_d}$  are i.i.d.  $\mathcal{B}(p)$

$H_1$ : There exists

$\mathcal{R}(i_1, i_2, \dots, i_d) = [i_1 - 1, i_1 + m_1 - 1] \times \dots \times [i_d - 1, i_d + m_d - 1] \subset \mathcal{R}_d$   
 where the r.v.'s  $X_{s_1, s_2, \dots, s_d} \sim \mathcal{B}(p')$ ,  $p' > p$  and  $X_{s_1, s_2, \dots, s_d} \sim \mathcal{B}(p)$   
 outside  $\mathcal{R}(i_1, i_2, \dots, i_d)$

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## Exact methods (Bernoulli case)

There are three main approaches used for investigating the exact distribution of the one dimensional discrete scan statistics over a sequence of binary trials:

- Combinatorial method: [Naus, 1974], [Naus, 1982]
- Finite Markov chain imbedding technique: [Fu, 2001], [Balakrishnan and Koutras, 2002], [Fu and Lou, 2003], [Wu, 2013], [Fu et al., 2012] etc.
- Conditional generating function method: [Ebnesahrashoob and Sobel, 1990], [Ebnesahrashoob et al., 2005], [Shinde and Kotwal, 2008] etc.

# Product Type Approximation

Considering that  $T_1 = L_1 m_1$ , [Naus, 1982] gave the following approximation

$$\mathbb{P}(S_{m_1}(T_1) \leq \tau) \approx Q(2m_1) \left[ \frac{Q(3m_1)}{Q(2m_1)} \right]^{\frac{T_1}{m_1} - 2},$$

where  $Q(2m_1) = \mathbb{P}(S_{m_1}(2m_1) \leq \tau)$  and  $Q(3m_1) = \mathbb{P}(S_{m_1}(3m_1) \leq \tau)$ .

$$Q(2m_1) = F^2(\tau; m_1, \rho) - \tau b(\tau + 1; m_1, \rho) F(\tau - 1; m_1, \rho) + m_1 \rho b(\tau + 1; m_1, \rho) F(\tau - 2; m_1 - 1),$$

$$Q(3m_1) = F^3(\tau; m_1, \rho) - A_1 + A_2 + A_3 - A_4,$$

$$A_1 = 2b(\tau + 1; m_1, \rho) F(\tau; m_1, \rho) [\tau F(\tau - 1; m_1, \rho) - m_1 \rho F(\tau - 2; m_1 - 1, \rho)],$$

$$A_2 = 0.5b^2(\tau + 1; m_1, \rho) [\tau(\tau - 1)F(\tau - 2; m_1, \rho) - 2(\tau - 1)m_1 F(\tau - 3; m_1 - 1, \rho) + m_1(m_1 - 1)\rho^2 F(\tau - 4; m_1 - 2, \rho)],$$

$$A_3 = \sum_{r=1}^{\tau} b(2(\tau + 1) - r; m_1, \rho) F^2(r - 1; m_1, \rho),$$

$$A_4 = \sum_{r=2}^{\tau} b(2(\tau + 1) - r; m_1, \rho) b(r + 1; m_1, \rho) [rF(r - 1; m_1, \rho) - m_1 \rho F(r - 2; m_1 - 1, \rho)].$$

# Product Type Approximation for Binomial and Poisson

If  $X_i \sim \text{Bin}(n, p)$  or  $X_i \sim \text{Pois}(\lambda)$ , we have the approximation

$$\begin{aligned} \mathbb{P}(S_{m_1}(T_1) \leq \tau) &\approx Q(2m_1) \left[ \frac{Q(3m_1)}{Q(2m_1)} \right]^{\frac{T_1}{m_1} - 2}, \quad T_1 \geq 3m_1 \\ &\approx Q(2m_1 - 1) \left[ \frac{Q(2m_1)}{Q(2m_1 - 1)} \right]^{T_1 - 2m_1 + 1}, \quad T_1 \geq 2m_1 \end{aligned}$$

where  $Q(2m_1 - 1)$ ,  $Q(2m_1)$  and  $Q(3m_1)$  are computed by [Karwe and Naus, 1997] recurrence.

► Karwe Naus algorithm for  $Q(2m_1 - 1)$  and  $Q(2m_1)$

## Bounds

[Glaz and Naus, 1991] developed a variety of tight bounds:

- Lower Bounds

$$\begin{aligned} \mathbb{P}(S_{m_1}(T_1) \leq \tau) &\leq \frac{Q(2m_1)}{\left[1 + \frac{Q(2m_1-1) - Q(2m_1)}{Q(2m_1-1)Q(2m_1)}\right]^{T_1 - 2m_1}}, \quad T_1 \geq 2m_1 \\ &\leq \frac{Q(3m_1)}{\left[1 + \frac{Q(2m_1-1) - Q(2m_1)}{Q(3m_1-1)}\right]^{T_1 - 3m_1}}, \quad T_1 \geq 3m_1 \end{aligned}$$

- Upper Bounds

$$\begin{aligned} \mathbb{P}(S_{m_1}(T_1) \leq \tau) &\leq Q(2m_1) [1 - Q(2m_1 - 1) + Q(2m_1)]^{T_1 - 2m_1}, \quad T_1 \geq 2m_1 \\ &\leq Q(3m_1) [1 - Q(2m_1 - 1) + Q(2m_1)]^{T_1 - 3m_1}, \quad T_1 \geq 3m_1 \end{aligned}$$

The values  $Q(2m_1 - 1)$ ,  $Q(2m_1)$ ,  $Q(3m_1 - 1)$ ,  $Q(3m_1)$  are computed using [Karwe and Naus, 1997] algorithm.

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# Product Type Approximation Bernoulli Case

[Boutsikas and Koutras, 2000] using Markov Chain Imbedding approach proposed the approximation

$$\mathbb{P}(S_{m_1, m_2}(T_1, T_2) \leq \tau) \approx \frac{Q(m_1, m_2)^{(T_1 - m_1 - 1)(T_2 - m_2 - 1)} Q(m_1 + 1, m_2 + 1)^{(T_1 - m_1)(T_2 - m_2)}}{Q(m_1, m_2 + 1)^{(T_1 - m_1 - 1)(T_2 - m_2)} Q(m_1 + 1, m_2)^{(T_1 - m_1)(T_2 - m_2 - 1)}}$$

Where,

$$Q(m_1, m_2) = F(\tau; m_1, m_2, \rho)$$

$$Q(m_1 + 1, m_2) = \sum_{s=0}^{\tau} F^2(\tau - s; m_2, \rho) b(s; (m_1 - 1)m_2, \rho)$$

$$Q(m_1 + 1, m_2 + 1) = \sum_{s_1, s_2=0}^{\tau} \sum_{t_1, t_2=0}^{\tau} \sum_{i_1, i_2, i_3, i_4=0}^1 b(s_1; m_1 - 1, \rho) b(s_2; m_1 - 1, \rho) b(t_1; m_2 - 1, \rho) \times$$

$$b(t_2; m_2 - 1, \rho) p^{\sum i_j} (1 - p)^{4 - \sum i_j} F(u; (m_1 - 1)(m_2 - 1), \rho)$$

$$u = \min \{ \tau - s_1 - t_1 - i_1, \tau - s_2 - t_1 - i_2, \tau - s_1 - t_2 - i_3, \tau - s_2 - t_2 - i_4 \}$$

$$b(s; n, \rho) = \binom{n}{s} \rho^s (1 - \rho)^{n-s}$$

$$F(s; n, \rho) = \sum_{i=0}^s b(i; n, \rho)$$

## Bounds for the Bernoulli Case

The following bounds were established by [Boutsikas and Koutras, 2003]

- Lower Bound

$$LB = (1 - Q_1)^{(T_1 - m_1)(T_2 - m_2)} (1 - Q_2)^{T_1 - m_1} (1 - Q_3)^{T_2 - m_2} (1 - Q_4)$$

- Upper Bound

$$UB = (1 - Q_1) \left(1 - q^{(m_1 - 1)(3m_2 - 2) + (2m_1 - 1)(m_2 - 1)} Q_1\right)^{(T_1 - m_1 - 1)(T_2 - m_2 - 1)} \left(1 - q^{m_1(m_2 - 1)} Q_1\right)^{T_2 - m_2 - 1} \\ \times \left(1 - q^{(m_1 - 1)(2m_2 - 1) + (m_1 - 1)(m_2 - 1)} Q_1\right)^{T_1 - m_1 - 1} \left(1 - q^{(m_1 - 1)(2m_2 - 1) + m_1(m_2 - 1)} Q_2\right)^{T_1 - m_1} \\ \times \left(1 - q^{(m_1 - 1)(3m_2 - 2) + m_1(m_2 - 1) + (m_1 - 1)(m_2 - 1)} Q_3\right)^{T_2 - m_2} \left(1 - q^{(m_1 - 1)(2m_2 - 1) + m_1(m_2 - 1)} Q_4\right).$$

Where  $X_{ij} \sim B(p)$ ,  $q = 1 - p$  and

$$Q_1 = F_{\tau+1, m_1 m_2}^c - q^{m_2} F_{\tau+1, (m_1 - 1)m_2}^c - q^{m_1} F_{\tau+1, m_1(m_2 - 1)}^c + q^{m_1 + m_2 - 1} F_{\tau+1, (m_1 - 1)(m_2 - 1)}^c, \\ Q_2 = F_{\tau+1, m_1 m_2}^c - q^{m_2} F_{\tau+1, (m_1 - 1)m_2}^c, \quad Q_3 = F_{\tau+1, m_1 m_2}^c - q^{m_1} F_{\tau+1, m_1(m_2 - 1)}^c, \quad Q_4 = F_{\tau+1, m_1 m_2}^c \\ F_{i, m}^c = 1 - F(i - 1; m, p).$$

# Product Type Approximation Binomial and Poisson

For  $X_{ij} \sim \text{Bin}(n, p)$  or  $X_{ij} \sim \text{Pois}(\lambda)$ , [Chen and Glaz, 2009] proposed the product type approximation

$$\mathbb{P}(S_{m_1, m_2}(T_1, T_2) \leq \tau) \approx \frac{Q(m_1+1, m_2+1)^{(T_1-m_1)(T_2-m_2)}}{Q(m_1+1, m_2)^{(T_1-m_1)(T_2-m_2-1)}} \times \frac{Q(m_1, 2m_2-1)^{(T_1-m_1-1)(T_2-2m_2)}}{Q(m_1, 2m_2)^{(T_1-m_1-1)(T_2-2m_2+1)}}$$

Where,

$$\begin{aligned} Q(m_1, 2m_2 - 1) &= \mathbb{P}(S_{m_1, m_2}(m_1, 2m_2 - 1) \leq \tau) \\ Q(m_1, 2m_2) &= \mathbb{P}(S_{m_1, m_2}(m_1, 2m_2) \leq \tau) \end{aligned}$$

To compute the unknown variables we use

- $Q(m_1, 2m_2 - 1)$  and  $Q(m_1, 2m_2)$  - adaptation of [Karwe and Naus, 1997] algorithm
- $Q(m_1 + 1, m_2)$  and  $Q(m_1 + 1, m_2 + 1)$  - conditioning

▶ Formulas for  $Q(m_1 + 1, m_2)$  and  $Q(m_1 + 1, m_2 + 1)$



## Lower Bound for Binomial and Poisson

[Chen and Glaz, 1996] gave a lower bound applying Hoover Bonferroni type inequality of order  $r \geq 3$ ,

$$\begin{aligned} \mathbb{P}(S_{m_1, m_2}(T_1, T_2) \geq \tau) &= \mathbb{P}\left(\bigcup_{i_1=1}^{T_1-m_1+1} \bigcup_{i_2=1}^{T_2-m_2+1} A_{i_1, i_2}\right) \leq \sum_{i_1=1}^{T_1-m_1+1} \sum_{i_2=1}^{T_2-m_2+1} \mathbb{P}(A_{i_1, i_2}) \\ &\quad - \sum_{i_1=1}^{T_1-m_1+1} \sum_{i_2=1}^{T_2-m_2} \mathbb{P}(A_{i_1, i_2} \cap A_{i_1, i_2+1}) - \sum_{i_1=1}^{T_1-m_1} \mathbb{P}(A_{i_1, 1} \cap A_{i_1+1, 1}) \\ &\quad - \sum_{i_1=1}^{T_1-m_1+1} \sum_{l=2}^{r-1} \sum_{i_2=1}^{T_2-m_2+1-l} \mathbb{P}(A_{i_1, i_2} \cap A_{i_1, i_2+1}^c \cdots A_{i_1, i_2+l-1}^c \cap A_{i_1, i_2+l}) \end{aligned}$$

Where  $A_{i_1, i_2} = \{Y_{i_1, i_2} \geq \tau\}$  and for  $r = 4$ ,

- Lower Bound

$$\begin{aligned} \mathbb{P}(S_{m_1, m_2}(T_1, T_2) \leq \tau) &\geq (T_1 - m_1)(Q(m_1 + 1, m_2) - 2Q(m_1, m_2)) - (T_1 - m_1 + 1)(T_2 - m_2 - 3) \\ &\quad \times Q(m_1, m_2 + 2) + (T_1 - m_1 + 1)(T_2 - m_2 - 2)Q(m_1, m_2 + 3). \end{aligned}$$

- $Q(m_1 + 1, m_2)$ ,  $Q(m_1, m_2)$ ,  $Q(m_1, m_2 + 2)$ ,  $Q(m_1, m_2 + 3)$  - [Karwe and Naus, 1997] algorithm (variant)

# Upper Bound for Binomial and Poisson

For the upper bound we adapt the inequality of [Kuai et al., 2000] to the two dimensional framework:

$$\begin{aligned} \mathbb{P}(S_{m_1, m_2}(T_1, T_2) \leq \tau) &= 1 - \mathbb{P}\left(\bigcup_{i_1=1}^{T_1 - m_1 + 1} \bigcup_{i_2=1}^{T_2 - m_2 + 1} A_{i_1, i_2}\right) \\ &\leq 1 - \sum_{i_1=1}^{T_1 - m_1 + 1} \sum_{i_2=1}^{T_2 - m_2 + 1} \left[ \frac{\theta_{i_1, i_2} \mathbb{P}(A_{i_1, i_2})^2}{\Sigma(i_1, i_2) + (1 - \theta_{i_1, i_2}) \mathbb{P}(A_{i_1, i_2})} + \frac{(1 - \theta_{i_1, i_2}) \mathbb{P}(A_{i_1, i_2})^2}{\Sigma(i_1, i_2) - \theta_{i_1, i_2} \mathbb{P}(A_{i_1, i_2})} \right] \end{aligned}$$

where

$$\Sigma(i_1, i_2) = \sum_{j_1=1}^{T_1 - m_1 + 1} \sum_{j_2=1}^{T_2 - m_2 + 1} \mathbb{P}(A_{i_1, i_2} \cap A_{j_1, j_2}) \quad \text{and} \quad \theta_{i_1, i_2} = \frac{\Sigma(i_1, i_2)}{\mathbb{P}(A_{i_1, i_2})} - \left\lfloor \frac{\Sigma(i_1, i_2)}{\mathbb{P}(A_{i_1, i_2})} \right\rfloor.$$

We have

$$\begin{aligned} \mathbb{P}(A_{i_1, i_2} \cap A_{j_1, j_2}) &= \begin{cases} [1 - Q(m_1, m_2)]^2, & \text{if } |i_1 - j_1| \geq m_1 \text{ or } |i_2 - j_2| \geq m_2, \\ 1 - 2Q(m_1, m_2) + \mathbb{P}(Y_{i_1, i_2} \leq \tau, Y_{j_1, j_2} \leq \tau), & \text{otherwise} \end{cases} \\ \mathbb{P}(Y_{i_1, i_2} \leq \tau, Y_{j_1, j_2} \leq \tau) &= \sum_{k=0}^{\tau} \mathbb{P}(Z = k) \mathbb{P}(Y_{i_1, i_2} - Z \leq \tau - k)^2, \\ Z &= \sum_{s=i_1 \vee j_1}^{(i_1 + m_1 - 1) \wedge (j_1 + m_1 - 1)} \sum_{t=i_2 \vee j_2}^{(i_2 + m_2 - 1) \wedge (j_2 + m_2 - 1)} X_{s, t}. \end{aligned}$$

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## Product Type Approximation

Glaz et al. proposed in [Guerriero et al., 2010] the product type approximation

$$\mathbb{P}(S_{m_1, m_2, m_3}(T_1, T_2, T_3) \leq \tau) \approx \frac{Q(m_1 + 1, m_2 + 1, m_3 + 1)^{(T_1 - m_1)(T_2 - m_2)(T_3 - m_3)} Q(m_1 + 1, m_2, m_3)^{(T_1 - m_1)(T_2 - m_2 - 1)(T_3 - m_3 - 1)}}{Q(m_1, m_2, m_3)^{(T_1 - m_1 - 1)(T_2 - m_2 - 1)(T_3 - m_3 - 1)} Q(m_1 + 1, m_2 + 1, m_3)^{(T_1 - m_1)(T_2 - m_2)(T_3 - m_3 - 1)}} \times \frac{Q(m_1, m_2 + 1, m_3)^{(T_1 - m_1 - 1)(T_2 - m_2)(T_3 - m_3 - 1)} Q(m_1, m_2, m_3 + 1)^{(T_1 - m_1 - 1)(T_2 - m_2 - 1)(T_3 - m_3 - 1)}}{Q(m_1 + 1, m_2, m_3 + 1)^{(T_1 - m_1 - 1)(T_2 - m_2 - 1)(T_3 - m_3 - 1)} Q(m_1, m_2 + 1, m_3 + 1)^{(T_1 - m_1 - 1)(T_2 - m_2)(T_3 - m_3 - 1)}}$$

Where,

$$Q(N_1, N_2, N_3) = \mathbb{P}(S_{m_1, m_2, m_3}(N_1, N_2, N_3) \leq \tau)$$

- The approximation also works for binomial and Poisson distribution
- Three Poisson Type Approximation

Lower and upper bounds for the distribution of three dimensional scan statistics were proposed by [Akiba and Yamamoto, 2004].

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# Haiman Type Approximation: Key Idea

Haiman proposed in [Haiman, 2000] a different approach

## Main Observation

The scan statistic r.v. can be viewed as a maximum of a sequence of 1-dependent stationary r.v..

- The idea:
  - discrete and continuous one dimensional scan statistic: [Haiman, 2000], [Haiman, 2007]
  - discrete and continuous two dimensional scan statistic: [Haiman and Preda, 2002], [Haiman and Preda, 2006]
  - discrete three dimensional scan statistic: [Amărioarei and Preda, 2013]

# Expressing $S_m(\mathbf{T})$ as a maximum of a 1-dependent sequence

Let  $L_j = \frac{T_j}{m_{j-1}}$ ,  $j \in \{1, 2, \dots, d\}$ , be positive integers

- Define for each  $k_1 \in \{1, 2, \dots, L_1 - 1\}$  the random variables

$$Z_{k_1} = \max_{\substack{(k_1-1)(m_1-1)+1 \leq i_1 \leq k_1(m_1-1) \\ 1 \leq j \leq (L_j-1)(m_j-1) \\ j \in \{2, \dots, d\}}} Y_{i_1, i_2, \dots, i_d}$$

- $(Z_j)_j$  is 1-dependent and stationary

Example:  $d = 1$

$$\underbrace{X_1, X_2, \dots, X_{m_1-1}}_{Z_1} \overbrace{X_{m_1}, \dots, X_{2(m_1-1)}}^{Z_2} \underbrace{X_{2m_1-1}, \dots, X_{3(m_1-1)}}_{Z_3} \overbrace{X_{3m_1-2}, \dots, X_{4(m_1-1)}}$$

- Observe

$$S_m(\mathbf{T}) = \max_{1 \leq k_1 \leq L_1 - 1} Z_{k_1}$$

## The main tool

Let  $(Z_j)_{j \geq 1}$  be a strictly stationary 1-dependent sequence of r.v.'s and let  $q_m = q_m(x) = \mathbb{P}(\max(Z_1, \dots, Z_m) \leq x)$ , with  $x < \sup\{u | \mathbb{P}(Z_1 \leq u) < 1\}$ .

### Theorem [Amărioarei, 2012]

For  $x$  such that  $\mathbb{P}(Z_1 > x) = 1 - q_1 < 0.1$  and  $m > 3$  we have

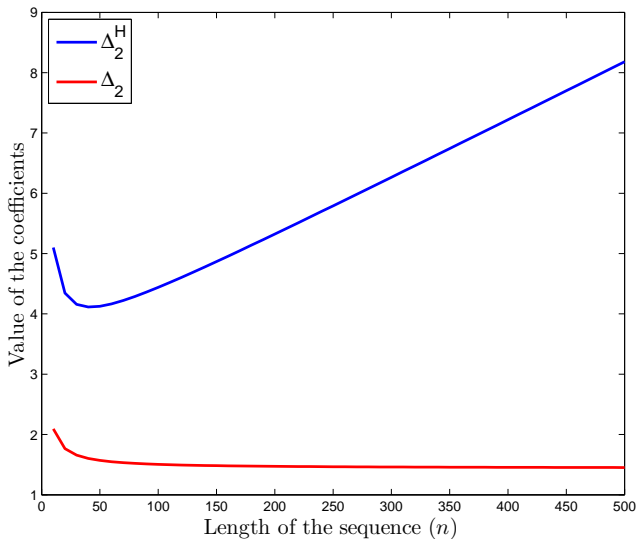
$$\left| q_m - \frac{2q_1 - q_2}{[1 + q_1 - q_2 + 2(q_1 - q_2)^2]^m} \right| \leq \Delta_2(1 - q_1)^2$$

- $\Delta_2 = mF(q_1, m) = m \left[ 1 + \frac{3}{m} + K(\alpha)(1 - q_1) + \frac{\Gamma(\alpha)(1 - q_1)}{m} \right]$ .

- Increased range of applicability
- Sharp bounds values (ex.  $\alpha = 0.025$ :  $561 \rightarrow 145$  and  $88 \rightarrow 17.5$ )



Difference between the results:  $1 - q_1 = 0.025$



## Approximation process

Define for  $t_1 \in \{2, 3\}$ ,

$$Q_{t_1} = Q_{t_1}(\tau) = \mathbb{P} \left( \bigcap_{k_1=1}^{t_1-1} \{Z_{k_1} \leq \tau\} \right) = \mathbb{P} \left( \begin{array}{l} 1 \leq i_1 \leq \max_{1 \leq j \leq (L_j-1)(m_j-1)} Y_{i_1, i_2, \dots, i_d} \leq \tau \\ 1 \leq i_j \leq (L_j-1)(m_j-1) \\ j \in \{2, \dots, d\} \end{array} \right)$$

If  $1 - Q_2 \leq 0.1$  then

$$\left| Q_m(T) - \frac{2Q_2 - Q_3}{[1 + Q_2 - Q_3 + 2(Q_2 - Q_3)^2]^{L_1 - 1}} \right| \leq (L_1 - 1)F(Q_2, L_1 - 1)(1 - Q_2)^2$$

Example:  $d = 1$

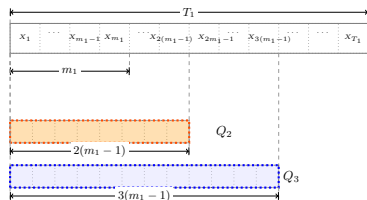
- The approximation

$$\mathbb{P}(S_{m_1}(T_1) \leq \tau) \approx H(Q_2, Q_3, L_1)$$

$$H(x, y, m) = \frac{2x - y}{[1 + x - y + 2(x - y)^2]^{m-1}}$$

- Approximation error, about

$$(L_1 - 1)F(Q_2, L_1 - 1)(1 - Q_2)^2$$



## Approximation process 2

The approximation of  $S_m(\mathbf{T})$  is an iterative process. The  $s$  step,  $1 \leq s \leq d$ , becomes:

- Let

$$Q_{t_1, t_2, \dots, t_s} = Q_{t_1, t_2, \dots, t_s}(\tau) = \mathbb{P} \left( \begin{array}{l} \max_{\substack{1 \leq i_l \leq (t_l - 1)(m_l - 1) \\ l \in \{1, \dots, s\}}} Y_{i_1, i_2, \dots, i_d} \leq \tau \\ 1 \leq i_j \leq (L_j - 1)(m_j - 1) \\ j \in \{s+1, \dots, d\} \end{array} \right)$$

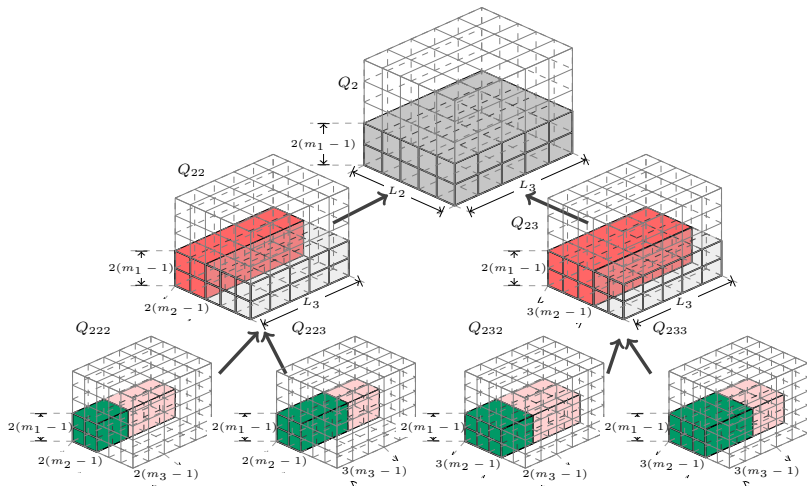
- Define for  $t_l \in \{2, 3\}$ ,  $l \in \{1, \dots, s-1\}$  and  $k_s \in \{1, 2, \dots, L_s - 1\}$

$$Z_{k_s}^{(t_1, t_2, \dots, t_{s-1})} = \begin{array}{l} \max_{\substack{1 \leq i_l \leq (t_l - 1)(m_l - 1) \\ l \in \{1, 2, \dots, s-1\}}} Y_{i_1, i_2, \dots, i_d} \\ (k_s - 1)(m_s - 1) + 1 \leq i_s \leq k_s(m_s - 1) \\ 1 \leq i_j \leq (L_j - 1)(m_j - 1) \\ j \in \{s+1, \dots, d\} \end{array}$$

- $\left\{ Z_1^{(t_1, t_2, \dots, t_{s-1})}, \dots, Z_{L_s - 1}^{(t_1, t_2, \dots, t_{s-1})} \right\}$  forms a 1-dependent stationary sequence
- The approximation

$$\left| Q_{t_1, \dots, t_{s-1}} - H \left( Q_{t_1, \dots, t_{s-1}, 2}, Q_{t_1, \dots, t_{s-1}, 3}, L_s \right) \right| \leq (L_s - 1) F(Q_{t_1, \dots, t_{s-1}, 2}, L_s - 1) (1 - Q_{t_1, \dots, t_{s-1}, 2})^2$$

# Illustration of the approximation process: $d = 3$



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Comparison between methods:  $d = 1$ Table 1 :  $n = 1, p = 0.005, m_1 = 10, T_1 = 1000, lt_{App} = 10^4$ 

$\tau$	Exact	Glaz and Naus Product type	Haiman Approximation	Approximation Error	Lower Bound	Upper Bound
1	0.810209	0.810216	0.810404	0.001111	0.809903	0.810439
2	0.995764	0.995764	0.995764	$3 \times 10^{-7}$	0.995764	0.995764
3	0.999950	0.999950	0.999950	$4 \times 10^{-11}$	0.999950	0.999950

Table 2 :  $n = 5, p = 0.05, m_1 = 25, T_1 = 500, lt_{App} = 10^4, lt_{Sim} = 10^3$ 

$\tau$	$\hat{\mathbb{P}}(S \leq \tau)$	Glaz and Naus Product type	Haiman Approximation	Total Error	Lower Bound	Upper Bound
13	0.712750	0.705787	0.714699	0.039308	0.697431	0.706948
14	0.867498	0.862184	0.865029	0.012502	0.859543	0.862407
15	0.946912	0.943329	0.946177	0.004169	0.942552	0.943362
16	0.980230	0.978959	0.979822	0.001354	0.978733	0.978963
17	0.993486	0.992821	0.993134	0.000433	0.992756	0.992822
18	0.997802	0.997726	0.997849	0.000127	0.997708	0.997726
19	0.999362	0.999327	0.999358	$3 \times 10^{-5}$	0.999322	0.999327
20	0.999819	0.999813	0.999825	$9 \times 10^{-6}$	0.999812	0.999813
21	0.999954	0.999951	0.999953	$2 \times 10^{-6}$	0.999951	0.999951

Comparison between methods:  $d = 2$ Table 3 :  $n = 1, p = 0.005, m_1 = m_2 = 6, T_1 = T_2 = 30, lt_{App} = 10^3, lt_{Sim} = 10^3$ 

$\tau$	$\hat{\mathbb{P}}(S \leq \tau)$	Glaz and Naus Product type	Haiman Approximation	Total Error(App+Sim)	Lower Bound	Upper Bound
2	0.915903	0.914013	0.920211	0.041483	0.901935	0.945623
3	0.994292	0.994395	0.994578	0.000803	0.993785	0.996638
4	0.999747	0.999757	0.999760	$2 \times 10^{-5}$	0.999737	0.999858
5	0.999992	0.999992	0.999992	$7 \times 10^{-7}$	0.999992	0.999995

Table 4 :  $n = 5, p = 0.002, m_1 = 5, m_2 = 10, T_1 = 50, T_2 = 80, lt_{App} = 10^4, lt_{Sim} = 10^3$ 

$\tau$	$\hat{\mathbb{P}}(S \leq \tau)$	Glaz and Naus Product type	Haiman Approximation	Total Error(App+Sim)	Lower Bound	Upper Bound
4	0.894654	0.873256	0.893724	0.037136	0.803422	0.944318
5	0.988003	0.986249	0.988144	0.002125	0.981418	0.993451
6	0.998963	0.998847	0.998963	0.000152	0.998543	0.999401
7	0.999926	0.999919	0.999925	$9 \times 10^{-6}$	0.999903	0.999955
8	0.999995	0.999995	0.999995	$5 \times 10^{-7}$	0.999994	0.999997

Comparison between methods:  $d = 3$ Table 5 :  $n = 1, p = 0.00005, m_1 = m_2 = m_3 = 5, T_1 = T_2 = T_3 = 60, lt_{App} = 10^5$ 

$\tau$	$\hat{\mathbb{P}}(S \leq \tau)$	Glaz et al. Product type	Haiman Approximation	Approximation Error	Simulation Error	Total Error
1	0.841806	0.841424	0.851076	0.011849	0.064889	0.076738
2	0.999119	0.999142	0.999192	0.000000	0.000170	0.000170
3	0.999997	0.999998	0.999997	0.000000	$3 \times 10^{-7}$	$3 \times 10^{-7}$

Table 6 :  $n = 1, p = 0.0001, m_1 = m_2 = m_3 = 5, T_1 = T_2 = T_3 = 60, lt_{App} = 10^5$ 

$\tau$	$\hat{\mathbb{P}}(S \leq \tau)$	Glaz et al. Product type	Haiman Approximation	Approximation Error	Simulation Error	Total Error
2	0.993294	0.993241	0.993192	0.000010	0.001367	0.001377
3	0.999963	0.999964	0.999963	0.000000	0.000005	0.000005
4	0.999999	0.999999	0.999999	0.000000	$2 \times 10^{-9}$	$2 \times 10^{-9}$



# Outline

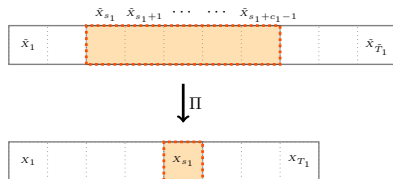
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# Introducing the model $d = 1$

Let  $m_1, c_1, \tilde{T}_1$  be positive integers such that  $c_1 \geq 1$  and  $2 \leq m_1 \leq T_1$

- Consider  $(\tilde{X}_i)_{1 \leq i \leq \tilde{T}_1}$  be a sequence of i.i.d. r.v.'s
- Take  $\Pi : \mathbb{R}^{c_1} \rightarrow \mathbb{R}$  to be a measurable function
- Define  $T_1 = \tilde{T}_1 - c_1 + 1$  and for  $1 \leq s_1 \leq T_1$

$$X_{s_1} = \Pi(\tilde{X}_{s_1}, \dots, \tilde{X}_{s_1+c_1-1})$$



- The sequence  $(X_{s_1})_{1 \leq s_1 \leq T_1}$  is  $(c_1 - 1)$ -dependent

## Approximation and bounds

Assume that  $L_1 = \frac{\tilde{T}_1}{m_1 + c_1 - 2}$  is a positive integer

- Define for  $k_1 \in \{1, 2, \dots, L_1 - 1\}$

$$Z_{k_1} = \max_{(k_1 - 1)(m_1 + c_1 - 2) + 1 \leq i_1 \leq k_1(m_1 + c_1 - 2)} Y_{i_1}$$

- $(Z_j)_j$  is 1-dependent and stationary
- We have the approximation

$$\mathbb{P}(S_{m_1}(T_1) \leq \tau) \approx \frac{2Q_2 - Q_3}{[1 + Q_2 - Q_3 + 2(Q_2 - Q_3)^2]^{L_1 - 1}}$$

- Error bounds

$$\begin{aligned} E_{app}(1) &= (L_1 - 1)F(Q_2, L_1 - 1)(1 - Q_2)^2, \\ E_{sapp}(1) &= (L_1 - 1)F(\hat{Q}_2, L_1 - 1)(1 - \hat{Q}_2 + \beta_2)^2, \\ E_{sf}(1) &= (L_1 - 1)(\beta_2 + \beta_3) \end{aligned}$$

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## Longest increasing run

Let  $\tilde{X}_1, \tilde{X}_2, \dots, \tilde{X}_{\tilde{T}_1}$  be a sequence of i.i.d. r.v.'s with the common distribution  $G$ .

A subsequence  $(\tilde{X}_k, \dots, \tilde{X}_{k+l-1})$  forms an *increasing run* (or ascending run) of length  $l \geq 1$ , starting at position  $k \geq 1$ , if it verifies the following relation

$$\tilde{X}_{k-1} > \tilde{X}_k < \tilde{X}_{k+1} < \dots < \tilde{X}_{k+l-1} > \tilde{X}_{k+l}$$

Denote the length of the longest increasing run among the first  $\tilde{T}_1$  random variables by  $M_{\tilde{T}_1}$ .

The asymptotic distribution was studied

- $G$  continuous distribution: [Pittel, 1981], [Révész, 1983], [Grill, 1987], [Novak, 1992], etc.
- $G$  discrete distribution: [Csaki and Foldes, 1996], [Grabner et al., 2003], [Eryilmaz, 2006], etc.

## Longest increasing run 2

Let  $c_1 = 2$ ,  $T_1 = \tilde{T}_1 - 1$  and define  $\Pi : \mathbb{R}^2 \rightarrow \mathbb{R}$  by

$$\Pi(x, y) = \begin{cases} 1, & \text{if } x < y \\ 0, & \text{otherwise} \end{cases}$$

Thus our block-factor model becomes

$$X_{s_1} = \mathbf{1}_{\tilde{X}_{s_1} < \tilde{X}_{s_1+1}}$$

Let  $L_{T_1}$  be the distribution of the length of the longest run of ones, among the first  $T_1$  observations.

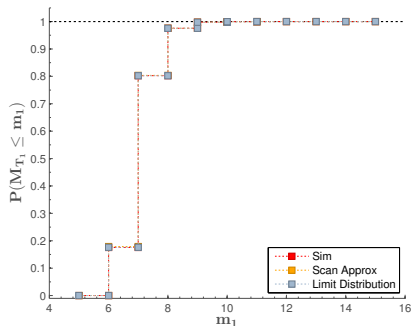
$$\mathbb{P}(M_{\tilde{T}_1} \leq m_1) = \mathbb{P}(L_{T_1} < m_1) = \mathbb{P}(S_{m_1}(T_1) < m_1), \text{ for } m_1 \geq 1$$

# Longest increasing run: numerical results

For  $\tilde{X}_{s_1} \sim \mathcal{U}([0, 1])$ , [Novak, 1992] showed that

$$\max_{1 \leq m_1 \leq T_1} \left| \mathbb{P}(L_{T_1} < m_1) - e^{-T_1 \frac{m_1+1}{(m_1+2)!}} \right| = \mathcal{O}\left(\frac{\ln T_1}{T_1}\right)$$

$m_1$	Sim	AppH	$E_{total}(1)$	LimApp
5	0.00000700	0.00000733	0.14860299	0.00000676
6	0.17567262	0.17937645	0.01089628	0.17620431
7	0.80257424	0.80362353	0.00110990	0.80215088
8	0.97548510	0.97566460	0.00011579	0.97550345
9	0.99749821	0.99751049	0.00001114	0.99749792
10	0.99977074	0.99977183	0.00000098	0.99977038
11	0.99998075	0.99998083	0.00000008	0.99998073
12	0.99999851	0.99999851	0.00000001	0.99999851
13	0.99999989	0.99999989	0.00000000	0.99999989
14	0.99999999	0.99999999	0.00000000	0.99999999
15	1.00000000	1.00000000	0.00000000	1.00000000



## Longest increasing run: numerical results

For  $\tilde{X}_{S_1} \sim \text{Geom}(p)$ , [Louchard and Prodinger, 2003] showed that

$$\mathbb{P}(M_{T_1} \leq m_1) \sim \exp(-\exp \eta),$$





$$\eta = \frac{m_1(m_1 + 1)}{2} \log \frac{1}{1-p} + m_1 \log \frac{1}{p} - \log T_1 - \log p + \log D(m_1),$$






$$D(m_1) = \prod_{k=1}^{m_1} [1 - (1-p)^k] [1 - (1-p)^{m_1+2}]$$






$m_1$	Sim	AppH	$E_{total}(1)$	LimApp
6	0.56445934	0.56997462	0.00255592	0.56810748
7	0.95295406	0.95325180	0.00018554	0.95294598
8	0.99658057	0.99659071	0.00001214	0.99657969
9	0.99979460	0.99979550	0.00000068	0.99979435
10	0.99998950	0.99998950	0.00000003	0.99998947





We used  $T_1 = 10000$ ,  $p = 0.1$  and  $lter = 10^5$ .








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# Computation of $Q(m_1 + 1, m_2)$ and $Q(m_1 + 1, m_2 + 1)$

Consider  $X_{ij} \sim B(n, p)$  and the notation  $Y_{i_1, i_2}^{j_1, j_2} = \sum_{i_1=1}^{j_1} \sum_{i_2=1}^{j_2} X_{ij}$ ,

$$\mathbb{P}(S_{m_1, m_2}(m_1 + 1, m_2) \leq k) = \sum_{y=0}^{k \wedge (m_1 - 1)m_2 n} \mathbb{P}^2(Y_{1,1}^{1, m_2} \leq k - y) \mathbb{P}(Y_{2,1}^{m_1, m_2} = y)$$

$$\begin{aligned} \mathbb{P}(S_{m_1, m_2}(m_1 + 1, m_2 + 1) \leq k) &= \sum_{y_1=0}^{k \wedge (m_1 - 1)(m_2 - 1)n} \sum_{y_2=0}^{(k - y_1) \wedge (m_2 - 1)n} \sum_{y_3=0}^{(k - y_1) \wedge (m_2 - 1)n} \\ &\quad \sum_{y_4=0}^{[k - y_1 - y_2 \vee y_3] \wedge (m_1 - 1)n} \sum_{y_5=0}^{[k - y_1 - y_2 \vee y_3] \wedge (m_1 - 1)n} \mathbb{P}(Y_{1,1}^{1,1} \leq a_1) \\ &\quad \times \mathbb{P}(Y_{1, m_2 + 1}^{1, m_2 + 1} \leq a_2) \mathbb{P}(Y_{m_1 + 1, 1}^{m_1 + 1, 1} \leq a_3) \mathbb{P}(Y_{m_1 + 1, m_2 + 1}^{m_1 + 1, m_2 + 1} \leq a_4) \\ &\quad \times \mathbb{P}(Y_{2,2}^{m_1, m_2} = y_1) \mathbb{P}(Y_{1,2}^{1, m_2} = y_2) \mathbb{P}(Y_{m_1 + 1, 2}^{m_1 + 1, m_2} = y_3) \\ &\quad \times \mathbb{P}(Y_{2,1}^{m_1, 1} = y_4) \mathbb{P}(Y_{2, m_2 + 1}^{m_1, m_2 + 1} = y_5) \\ a_1 &= k - y_1 - y_2 - y_4, \quad a_2 = k - y_1 - y_2 - y_5, \\ a_3 &= k - y_1 - y_3 - y_4, \quad a_4 = k - y_1 - y_3 - y_5, \end{aligned}$$

# Karwe Naus recursive methods

Define

$$b_{2(m)}^k(y) = \mathbb{P}(S_m(2m) \leq k, Y_{m+1}(m) = y)$$

$$f(y) = \mathbb{P}(X_1 = y)$$

$$Q_{2m}^k = \mathbb{P}(S_m(2m) \leq k)$$

We have the recurrence relations

$$b_{2(1)}^k(y) = \left( \sum_{j=0}^k f(j) \right) f(y)$$

$$b_{2(m)}^k(y) = \sum_{\eta=0}^y \sum_{\nu=0}^{k-y+\eta} b_{2(m-1)}^{k-\nu}(y-\eta) f(\nu) f(\eta)$$

$$Q_{2m}^k = \sum_{y=0}^k b_{2(m)}^k(y)$$

$$Q_{2m-1}^k = \sum_{x=0}^k f(x) Q_{2(m-1)}^{k-x}$$

## Selected Values for $K(\alpha)$ and $\Gamma(\alpha)$

$\alpha$	$K(\alpha)$	$\Gamma(\alpha)$
0.1	38.63	480.69
0.05	21.28	180.53
0.025	17.56	145.20
0.01	15.92	131.43

Table 7 : Selected values for  $K(\alpha)$  and  $\Gamma(\alpha)$

← Return

# Error bounds

Let  $\gamma_{t_1, \dots, t_d} = Q_{t_1, \dots, t_d}$ , with  $t_j \in \{2, 3\}$ ,  $j \in \{1, \dots, d\}$ , and define for  $2 \leq s \leq d$

$$\gamma_{t_1, \dots, t_{s-1}} = H(\gamma_{t_1, \dots, t_{s-1}, 2}, \gamma_{t_1, \dots, t_{s-1}, 3}, L_s)$$

Denote with  $\hat{Q}_{t_1, \dots, t_d}$  the estimated value of  $Q_{t_1, \dots, t_d}$  and define for  $2 \leq s \leq d$

$$\hat{Q}_{t_1, \dots, t_{s-1}} = H(\hat{Q}_{t_1, \dots, t_{s-1}, 2}, \hat{Q}_{t_1, \dots, t_{s-1}, 3}, L_s)$$

We observe that

$$\left| Q_m(T) - H(\hat{Q}_2, \hat{Q}_3, L_1) \right| \leq \underbrace{\left| Q_m(T) - H(\gamma_2, \gamma_3, L_1) \right|}_{E_{app}(d)} + \underbrace{\left| H(\gamma_2, \gamma_3, L_1) - H(\hat{Q}_2, \hat{Q}_3, L_1) \right|}_{E_{sim}(d)}$$

Return

# Error bounds: approximation and simulation errors 1

## Approximation error

$$E_{app}(d) = \sum_{s=1}^d (L_1 - 1) \cdots (L_s - 1) \sum_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1} \in \{2,3\}} F_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}} \left( 1 - \gamma_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2} + B_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2} \right)^2,$$

where for  $2 \leq s \leq d$

$$F_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}} = F(Q_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2}, L_s - 1), \quad F = F(Q_2, L_1 - 1),$$

$$B_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}} = (L_s - 1) \left[ F_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}} \left( 1 - \gamma_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2} + B_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2} \right)^2 + \sum_{\mathbf{t}_s \in \{2,3\}} B_{\mathbf{t}_1, \dots, \mathbf{t}_s} \right],$$

$$B_{\mathbf{t}_1, \dots, \mathbf{t}_{d-1}} = (L_d - 1) F_{\mathbf{t}_1, \dots, \mathbf{t}_{d-1}} \left( 1 - \gamma_{\mathbf{t}_1, \dots, \mathbf{t}_{d-1}, 2} + B_{\mathbf{t}_1, \dots, \mathbf{t}_{d-1}, 2} \right)^2, \quad B_{\mathbf{t}_1, \dots, \mathbf{t}_d} = 0,$$

and for  $s = 1$ :  $\sum_{\mathbf{t}_1, \mathbf{t}_0 \in \{2,3\}} x = x$ ,  $F_{\mathbf{t}_1, \mathbf{t}_0} = F$ ,  $\gamma_{\mathbf{t}_1, \mathbf{t}_0, 2} = \gamma_2$  and  $B_{\mathbf{t}_1, \mathbf{t}_0, 2} = B_2$ .

# Error bounds: approximation and simulation errors 2

## Simulation errors

$$E_{sf}(d) = (L_1 - 1) \dots (L_d - 1) \sum_{\mathbf{t}_1, \dots, \mathbf{t}_d \in \{2,3\}} \beta_{\mathbf{t}_1, \dots, \mathbf{t}_d}$$

$$E_{sapp}(d) = \sum_{s=1}^d (L_1 - 1) \dots (L_s - 1) \sum_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1} \in \{2,3\}} F_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}} \left(1 - \hat{Q}_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2} + A_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2} + C_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2}\right)^2$$

where for  $2 \leq s \leq d$

$$A_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}} = (L_s - 1) \dots (L_d - 1) \sum_{\mathbf{t}_s, \dots, \mathbf{t}_d \in \{2,3\}} \beta_{\mathbf{t}_1, \dots, \mathbf{t}_d}, \quad A_{\mathbf{t}_1, \dots, \mathbf{t}_d} = \beta_{\mathbf{t}_1, \dots, \mathbf{t}_d}$$

$$C_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}} = (L_s - 1) \left[ F_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}} \left(1 - \hat{Q}_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2} + A_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2} + C_{\mathbf{t}_1, \dots, \mathbf{t}_{s-1}, 2}\right)^2 + \sum_{\mathbf{t}_s \in \{2,3\}} C_{\mathbf{t}_1, \dots, \mathbf{t}_s} \right]$$